**XJTLU Entrepreneur College (Taicang) Cover Sheet**

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| --- | --- | --- |
| Module code and Title | Database Development and Design (CPT201TC) | |
| School Title | School of AI and Advanced Computing | |
| Assignment Title | 001: Assessment Task 1 (CW) | |
| Submission Deadline | 17:00, 10th Dec (Friday) | |
| Final Word Count | NA | |
| If you agree to let the university use your work anonymously for teaching and learning purposes, please type **“yes”** here. | | **Yes** |

I certify that I have read and understood the University’s Policy for dealing with Plagiarism, Collusion and the Fabrication of Data (available on Learning Mall Online). With reference to this policy I certify that:

* My work does not contain any instances of plagiarism and/or collusion.  
  My work does not contain any fabricated data.

**By uploading my assignment onto Learning Mall Online, I formally declare that all of the above information is true to the best of my knowledge and belief.**

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| **Scoring – For Tutor Use** | | | | | | |
| **Student ID** | | | |  | | |
|  | | | | | | |
| **Stage of Marking** | | **Marker**  **Code** | **Learning Outcomes Achieved （F/P/M/D）**  **(please modify as appropriate)** | | | **Final**  **Score** |
| **A** | **B** | **C** |
| 1st Marker – red pen | |  |  |  |  |  |
| Moderation  – green pen | | **IM**  **Initials** | The original mark has been accepted by the moderator (please circle as appropriate): | | | Y / N |
|  | Data entry and score calculation have been checked by another tutor (please circle): | | | Y |
| 2nd Marker if needed – green pen | |  |  |  |  |  |
| **For Academic Office Use** | | | **Possible Academic Infringement (please tick as appropriate)** | | | |
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|  |  |  | **Category B** | |
| **Category C** | |
| **Category D** | |
| **Category E** | |

**Students**

The assignment must be submitted in a single pdf document via Learning Mall Online to the correct drop box. Only electronic submission is accepted and no hard copy submission.

All students must download their file and check that it is viewable after submission. Documents may become corrupted during the uploading process (e.g. due to slow internet connections). However, students themselves are responsible for submitting a functional and correct file for assessments.

**Database Development and Design (CPT201TC)**

**Assessment 001: Individual Coursework**

**Due:** Dec 10th, 2021 @ 17:00

**Weight:** 50%

**Maximum Marks:** 50

**Overview**

This assignment aims to gain experience in understanding the internal functionality of different database management systems, including RDBMS, XML, Object-relational databases, and gain experience in designing data warehouses to perform OLAP operations. The course work will be assessed for the following learning outcomes:

1. Identify and apply the principles underpinning transaction management within DBMS.
2. Demonstrate an understanding of advanced SQL topics.
3. Illustrate the issues related to Web technologies and DBMS and XML as a semi-structured data representation formalism.
4. Identify the principles underlying object-relational models.
5. State the main concepts in data warehousing and data mining.

**Assessment Tasks**

Your task is to answer every question by carefully reading the questions and guidelines for the system setup. Record your thoughts and assumptions, where necessary, while reporting your answers. There are the following five parts for this course work corresponding to five learning outcomes of this module:

1. Advanced SQL, Triggers, Indexing, and Query Optimization
2. Transaction Management
3. Querying XML Data
4. Object-Relational Database
5. Data Warehousing and OLAP

**Marking Criteria**

This coursework will be graded out of 50% of the total marks with a maximum of 50 marks. There are five (5) parts and ten (10) marks available for each part. Marks will be awarded as follows:

**0 mark:** Not attempted/wrong

**50% marks:** Partially attempted/correct

**100% marks:** Complete and correct

**1: Advanced SQL, Triggers, and Indexing**

**(10 Marks)**

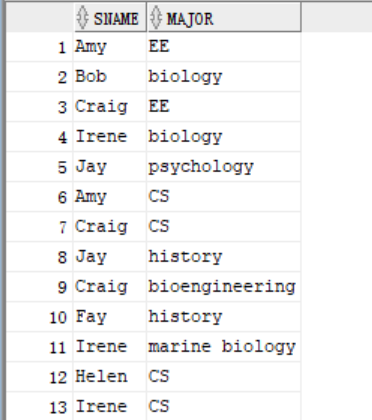
**Q1(a)** The dataset to be used in this exercise is the College\_Admission database provided along with this assignment. The dataset contains a normalized schema for a table of students, colleges, and information about students applying to colleges. You are required to load the data in Oracle Database and use SQL Developer to answer the following questions. For each question please provide query statements and results. **(1\*3=3 marks)**

1. Retrieve students’ names with majors for which they have applied.
2. Retrieve the name, GPA, and application decision of students from the high school with less than 1000 students and have applied to major in CS at Stanford.
3. Compute the amount by which the average GPA of students applying to CS exceeds the average GPA of students who are not applying to CS. [*Hint: compute two subqueries in the from clause, one of them computes the average GPA of CS applicants and one the average GPA of non CS applicants, in the select clause do the subtraction of the non-CS GPA from the CS-GPA*].

-- (a)i

select distinct SNAME, MAJOR

from student join apply on student.SID=apply.SID;



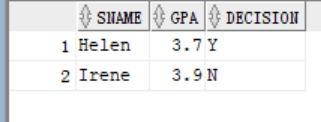
-- (a)ii

select SNAME,GPA, DECISION

from student join apply on student.SID=apply.SID

join college on college.CNAME=apply.CNAME

where SIZEHS<1000 and MAJOR='CS' and college.CNAME='Stanford';



-- (a)iii

select CS\_GPA-nonCS\_GPA

from

(

select avg(GPA) as CS\_GPA

from student join apply on student.SID=apply.SID

where MAJOR='CS'

) t1,

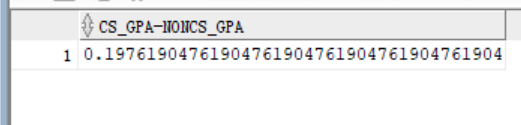
(

select avg(GPA) as nonCS\_GPA

from student join apply on student.SID=apply.SID

where MAJOR!='CS'

) t2;



**Q1(b)** In the next task, use the same database (College\_Admission) by using PL/SQL block to work on triggers for the scenario given below: **(4 marks)**

Create a row-level trigger on the student table that will fire after insertions into the student table by checking the GPA value of a new record. If the GPA of the inserted student is greater than 3.3, or less than or equal to 3.6, that student will be automatically applying to Stanford for a geology major and applying to MIT for a biology major. In other words, if the GPA of a new student satisfies the given condition (GPA>3.3 or GPA<=3.6), two new rows should be automatically added to Apply table, one with values (sID of the new Student, cName: ‘Stanford’, major: ‘geology’, Decision: NULL, and second row with values ( sID of the new student, cName: ‘MIT’, major: ‘biology’, Decision: NULL).

To test the above trigger, create and execute SQL statements to insert a new row in the student table with satisfying GPA conditions. E.g.

insert into Student values ('111', 'Kevin', 3.5, 1000);

insert into Student values ('222', 'Lori', 3.8, 1000);

create or replace trigger trg\_fire\_student

after insert on student for each row

begin

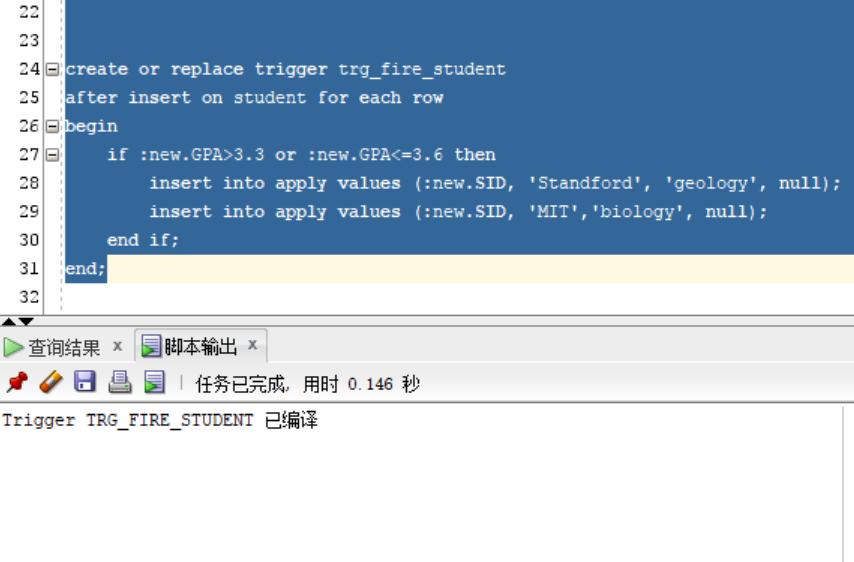
if :new.GPA>3.3 or :new.GPA<=3.6 then

insert into apply values (:new.SID, 'Standford', 'geology', null);

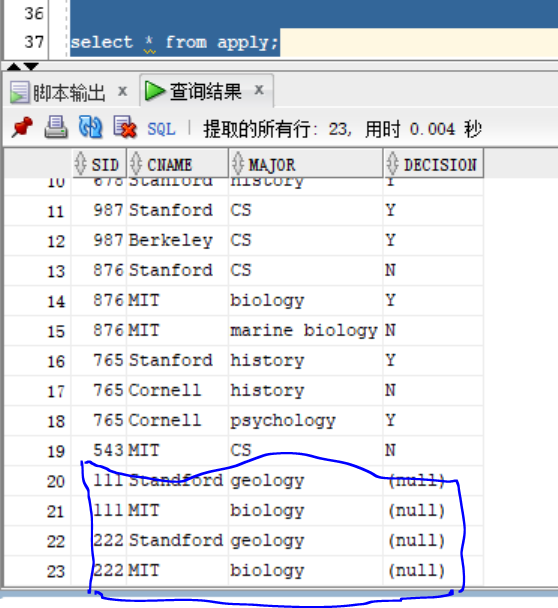
insert into apply values (:new.SID, 'MIT','biology', null);

end if;

end;



The first insert statement satisfies the GPA condition (GPA>3.3 or GPA<=3.6), so two new tuples should be automatically added in the Apply table. You should provide trigger code, test SQL and query results showing new records in the student, and Apply tables according to your insert statement.



**Q1(c)** In the next task, you will be working on the same database to work on indexes by considering the following scenario: **(3 marks)**

Create a tree-based index (one at a time) on the following attribute pairs (Note: by default, indexes are tree-based).

Student.sID, College.cName

Student.sID, Student.GPA

Apply.cName, College.cName

Apply.sID, Student.GPA

Execute the following query each time you create an index, and evaluate which two pairs of the index are most useful for speeding up query execution. For each pair, discuss why or why not that index pair is useful based on the number of disc blocks required to access to retrieve a record.

Select \* From Student, Apply, College

Where Student.sID = Apply.sID and Apply.cName = College.cName

And Student.GPA > 1.5 And College.cName < 'Cornell'

-- Student.sID, College.cName

create index idx\_sid on student(SID);

create index idx\_cname on college(CNAME);

-- Student.sID, Student.GPA

create index idx\_sid on student(SID);

create index idx\_gpa on student(GPA);

--Apply.cName, College.cName

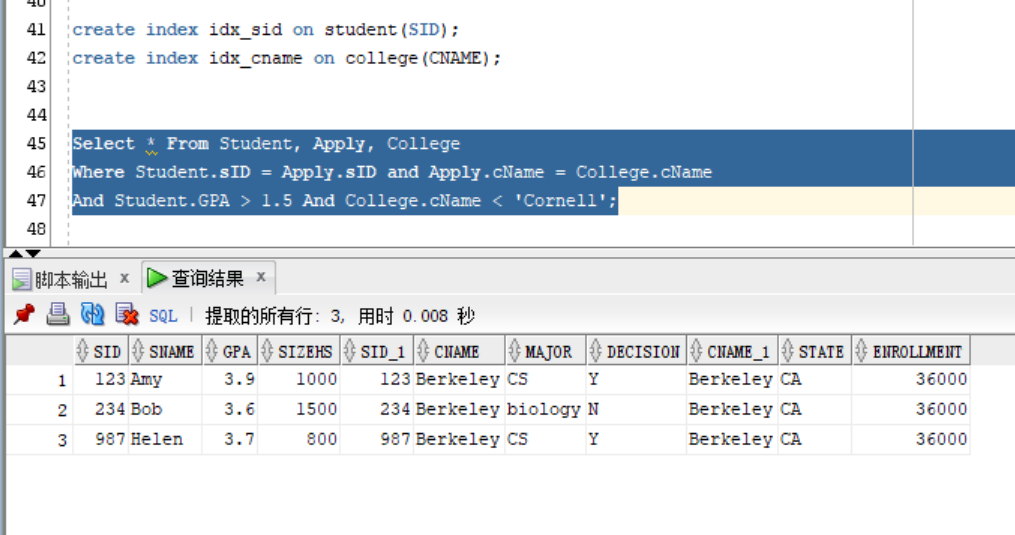
create index idx\_ap\_cname on apply(CNAME);

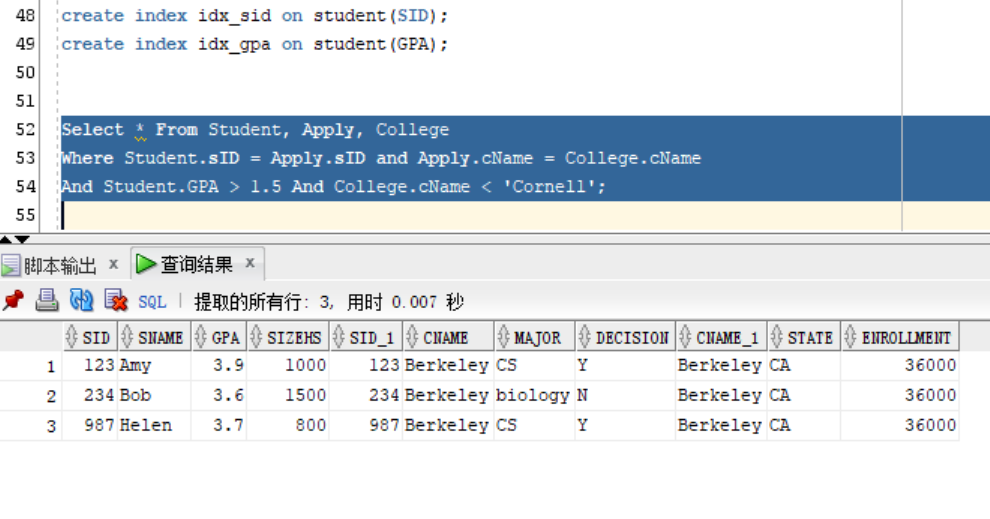
create index idx\_cname on college(CNAME);

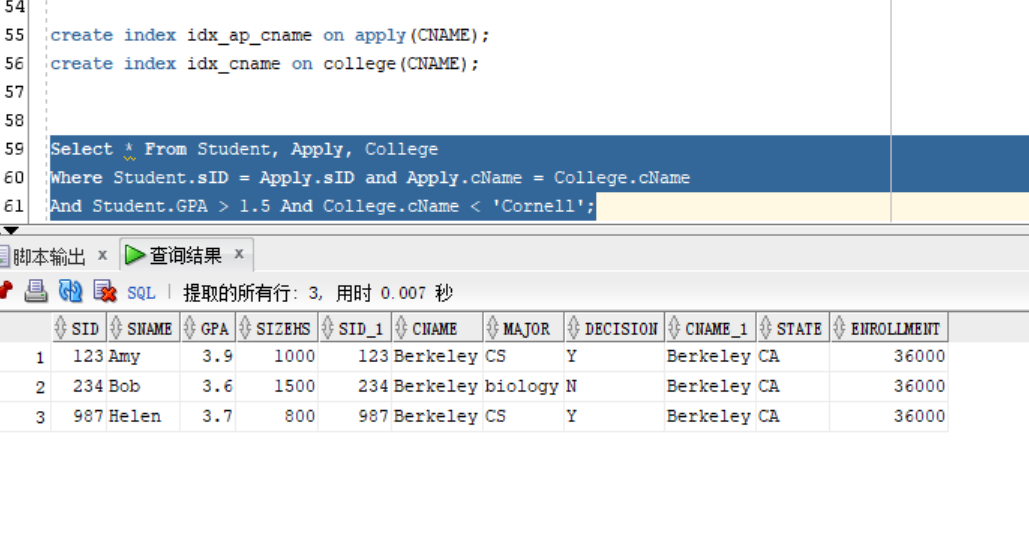
--Apply.sID, Student.GPA

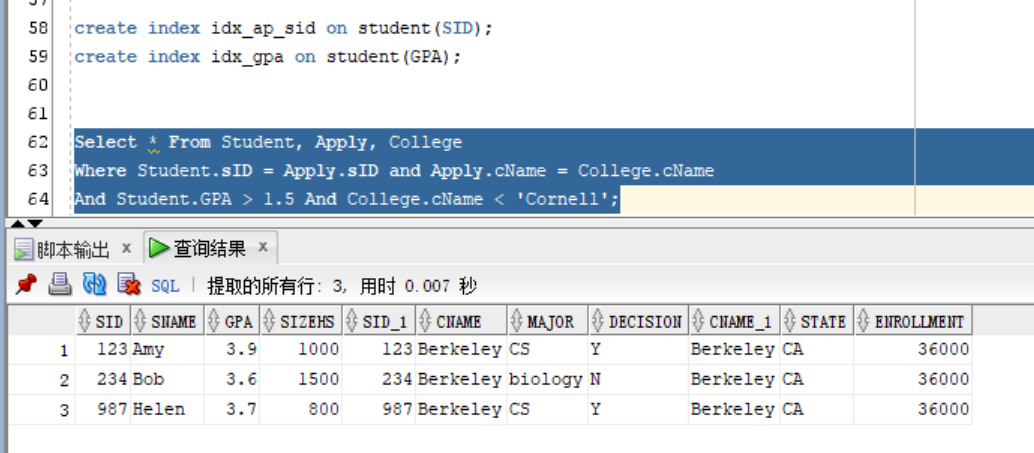
create index idx\_ap\_sid on student(SID);

create index idx\_gpa on student(GPA);









from all the query result, I can see the last three indexes groups used same query time, they all about 0.007s, and the first one is 0.008s. the last three all used on comparation condition but only the first one just create indexes on join conditions.

**2: Transaction Management (10 Marks)**

**Q2(a)**: Consider a relation R(A) containing {(5),(6)} and two transactions:

**T1**: Update R set A = A+1;

**T2**: Update R set A = 2\*A.

Suppose both transactions are submitted under the **isolation and atomicity** properties. Provide all possible final states of R (values of R) by considering all possible serializable and non-serializable schedules. **(2 marks)**

If both transactions are submitted under the isolation and atomicity properties.

a.1 all possible serializable schedules: T1 first then T2; T2 first and then T1

the beginning state of R is {(5),(6)}

T1 first then T2: final state of R is {(12),(14)}

|  |  |  |  |
| --- | --- | --- | --- |
| Time | T1 | T2 | A |
| t1 | read(A) |  | {(5),(6)} |
| t2 | write(A) |  | {(6),(7)} |
| t3 | commit |  | {(6),(7)} |
| t4 |  | read(A) | {(6),(7)} |
| t5 |  | write(A) | {(12),(14)} |
| t6 |  | commit | {(12),(14)} |

T2 first then T1: final state of R is {(11),(13)}

|  |  |  |  |
| --- | --- | --- | --- |
| Time | T1 | T2 | A |
| t1 |  | read(A) | {(5),(6)} |
| t2 |  | write(A) | {(10),(12)} |
| t3 |  | commit | {(10),(12)} |
| t4 | read(A) |  | {(10),(12)} |
| t5 | write(A) |  | {(11),(13)} |
| t6 | commit |  | {(11),(13)} |

a.2 all possible non-serializable schedules

the beginning state of R is {(5),(6)}

Firstly, T1 read, then T2 read, write and commit, then T1 write and commit, the final state of A is {(6),(7)}

|  |  |  |  |
| --- | --- | --- | --- |
| Time | T1 | T2 | A |
| t1 | read(A) |  | {(5),(6)} |
| t2 |  | read(A) | {(5),(6)} |
| t3 |  | write(A) | {(10),(12)} |
| t4 |  | commit | {(10),(12)} |
| t5 | write(A) |  | {(6),(7)} |
| t6 | commit |  | {(6),(7)} |

Firstly, T2 read, then T2 write, Then T1 read, and T2 commit, then T1 write and commit, the final state of A is {(6),(7)}

|  |  |  |  |
| --- | --- | --- | --- |
| Time | T1 | T2 | A |
| t1 |  | read(A) | {(5),(6)} |
| t2 |  | write(A) | {(10),(12)} |
| t3 | read(A) |  | {(5),(6)} |
| t4 |  | commit | {(10),(12)} |
| t5 | write(A) |  | {(6),(7)} |
| t6 | commit |  | {(6),(7)} |

Firstly, T2 read, Then T1 read, write and commit, and T2 write and commit, the final state of A is {(6),(7)}

|  |  |  |  |
| --- | --- | --- | --- |
| Time | T1 | T2 | A |
| t1 |  | read(A) | {(5),(6)} |
| t2 | read(A) |  | {(5),(6)} |
| t3 | write(A) |  | {(6),(7)} |
| t4 | commit |  | {(6),(7)} |
| t5 |  | write(A) | {(10),(12)} |
| t6 |  | commit | {(10),(12)} |

**Q2(b)** Consider a table R(A) containing {(1),(2)} and following two transactions run concurrently:

**T1**: Update R set A = 2\*A; commit;

**T2**: Select avg(A) from R; commit;

If transaction T2 executes using **READ UNCOMMITTED**, what are the possible average values of T2 returns? Provide an explanation for each value it returns. (2 marks)

If transaction T2 executes using read uncommitted.

b.1 if T1 commited firstly, then T2 read and compute, the result of avg(A) is 3

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 | write(A) |  | {(1),(2)} |  |
| t2 | commit; |  | {(2),(4)} |  |
| t3 |  | read(A) | {(2),(4)} | 3 |
| t4 |  | commit; | {(2),(4)} | 3 |

b.2 if T1 write first then T2 read and compute and commit, T1 committed finally, the result of avg(A) is 3

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 | write(A) |  | {(2),(4)} |  |
| t2 |  | read(A) | {(2),(4)} | 3 |
| t3 |  | commit; | {(2),(4)} | 3 |
| t4 | commit; |  | {(2),(4)} |  |

b.3 if T2 read first then T1 write and commit, T2 committed finally, the result of avg(A) is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 |  | read(A) | {(1),(2)} | 1.5 |
| t2 | write(A) |  | {(2),(4)} |  |
| t3 | commit; |  | {(2),(4)} |  |
| t4 |  | commit; | {(2),(4)} | 1.5 |

b.4 if T2 read first and commit, then T1 write and commit, the result of avg(A) is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 |  | read(A) | {(1),(2)} | 1.5 |
| t2 |  | commit; | {(1),(2)} | 1.5 |
| t3 | write(A) |  | {(2),(4)} |  |
| t4 | commit; |  | {(2),(4)} |  |

**Q2(c)** Consider table R(A) containing {(1),(2)}. Suppose we have the following two transactions running concurrently:

**T1:** Update T set A=2\*A;

insert into R values (6);

Commit;

**T2:** Select avg(A) from R;

Select avg(A) from R ;

Commit;

If transaction T2 executes using **REPEATABLE READ,** what are the possible values returned by T2 in its SECOND select statement? Provide an explanation for each value it returns. **(2 marks)**

If transaction T2 executes using REPEATABLE READ.

And the initial value of R(A) is {(1),(2)}

c.1 T1 first then T2, the result of T2’s select statement is 4.

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 | write(A) |  | {(1),(2)} |  |
| t2 | commit |  | {(2),(4),(6)} |  |
| t3 |  | read(A) | {(2),(4),(6)} | 4 |
| t4 |  | read(A) | {(2),(4),(6)} | 4 |
| t5 |  | commit |  |  |

c.2 T2 first schedule then T1, the result of T2’s select statement is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 |  | read(A) | {(1),(2)} | 1.5 |
| t2 |  | read(A) | {(1),(2)} | 1.5 |
| t3 |  | commit | {(1),(2)} |  |
| t4 | write(A) |  | {(1),(2)} |  |
| t5 | commit |  | {(2),(4),(6)} |  |

c.3 T1 wirtes first but not commit, then T2 read and commit, the result of T2’s select statement is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 | write(A) |  | {(1),(2)} |  |
| t2 |  | read(A) | {(1),(2)} | 1.5 |
| t3 |  | read(A) | {(1),(2)} | 1.5 |
| t4 |  | commit | {(1),(2)} |  |
| t5 | commit |  | {(2),(4),(6)} |  |

c.4 T1 wirtes first but not commit, then T2 read, then T1 commit, then T2 read and commit,the result of T2’s select statement is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 | write(A) |  | {(1),(2)} |  |
| t2 |  | read(A) | {(1),(2)} | 1.5 |
| t3 | commit |  | {(2),(4),(6)} | 1.5 |
| t4 |  | read(A) | {(2),(4),(6)} | 1.5 |
| t5 |  | commit | {(2),(4),(6)} | 1.5 |

c.5 T2 read first then T1 write and commit, then T2 read and commit,the result of T2’s select statement is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 |  | read(A) | {(1),(2)} | 1.5 |
| t2 | write(A) |  | {(1),(2)} | 1.5 |
| t3 | commit |  | {(2),(4),(6)} | 1.5 |
| t4 |  | read(A) | {(2),(4),(6)} | 1.5 |
| t5 |  | commit | {(2),(4),(6)} | 1.5 |

c.6 T2 read first then T1 write then T2 read , then T1 commit and T2 commit,the result of T2’s select statement is 1.5

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| time | T1 | T2 | A | avg(A) |
| t1 |  | read(A) | {(1),(2)} | 1.5 |
| t2 | write(A) |  | {(1),(2)} | 1.5 |
| t3 |  | read(A) | {(1),(2)} | 1.5 |
| t4 | commit |  | {(2),(4),(6)} | 1.5 |
| t5 |  | commit | {(2),(4),(6)} | 1.5 |

**Q2(d)** For each of the following schedules of read, write, commit and abort actions done by transactions T1, T2, T3, state whether they are recoverable or not. If not recoverable, what type of inconsistencies might we have in the final database state? If recoverable, which of the other transactions need to be rolled back? Note ri (A) and wi (A) mean that transaction Ti reads and writes database object A, respectively. **(2\*2=4 marks)**

|  |  |  |
| --- | --- | --- |
| **S1** | | |
| T1 | T2 | T3 |
| w1(A);  w1(B);  abort; | r2(A);  w2(A\*2)  r2(B);  commit; | r3(B);  r3(B+5);  commit; |

d.1 S1 can not recover. Because T2 read the item B which written by T1 and T1 hadn’t committed yet.

The final database may lost update. Because T1’s written on item B may be lost if T1 is abort.

|  |  |  |
| --- | --- | --- |
| **S2** | | |
| T1 | T2 | T3 |
| r1(A);  r1(B);  w1(B);  abort; | r2(A);  w2(A\*2) | r3(B);  w3(B+5); |

d.2 S2 can recover. T3 read the item B which written by T1 and not committed, and T3 also not committed, so it not violates the recoverable regulation.

If want to recover, T3, T1 need to roll back.

**3: Querying XML Data (10 Marks)**

**Q3** In this task, you are required to use the Company XML schema provided along with the lab manual (Chapter 3). Specify the following views by writing and executing XQuery expressions on the company.xml document. For each query, provide query expression and results. **(2\*5=10 marks)**

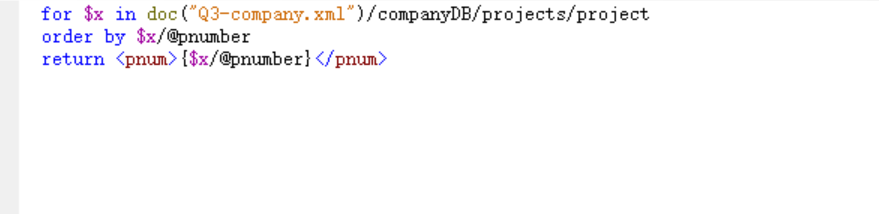
*Note: You can use any XML editor using the software guide provided along with the lab manual.*

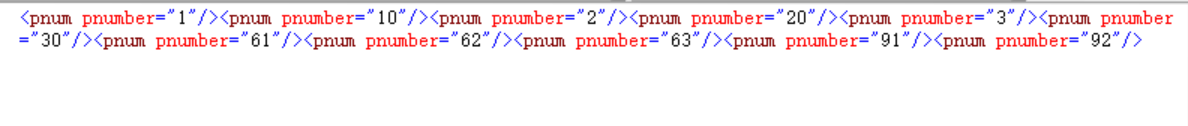
1. Get distinct project numbers of projects in which employees work. The results should be ordered by project number.

for $x in doc("Q3-company.xml")/companyDB/projects/project

order by $x/@pnumber

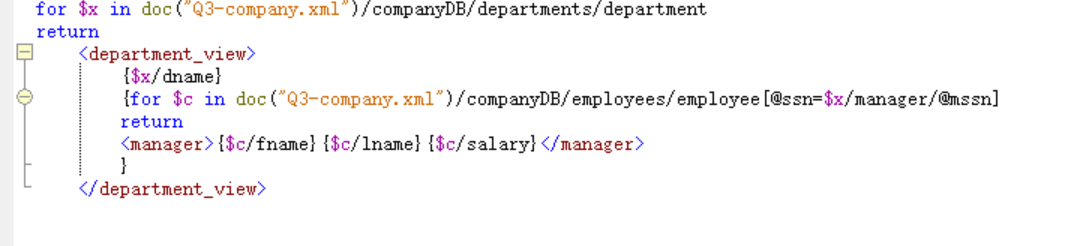
return <pnum>{$x/@pnumber}</pnum>

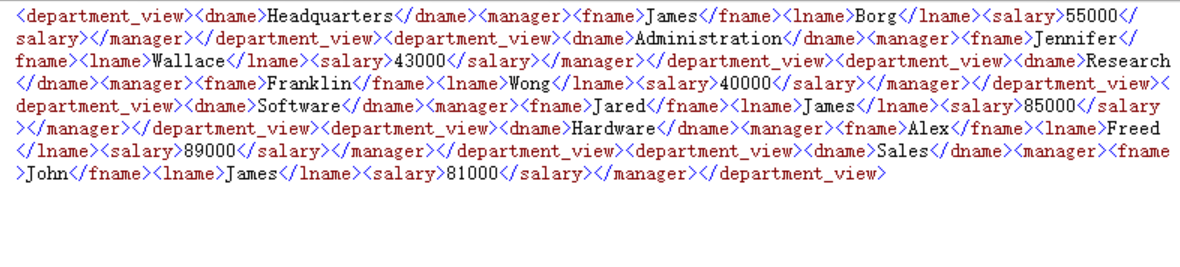




1. Create a view that has the department name, manager name, and manager salary for every department.

for $x in doc("Q3-company.xml")/companyDB/departments/department  
return  
 <department\_view>  
 {$x/dname}  
 {for $c in doc("Q3-company.xml")/companyDB/employees/employee[@ssn=$x/manager/@mssn]  
 return   
 <manager>{$c/fname}{$c/lname}{$c/salary}</manager>  
 }  
 </department\_view>





1. A view that has the employee name, supervisor name, and employee salary for each employee who works in the ‘Research’ department.

for $e in doc("Q3-company.xml")/companyDB/employees/employee  
return   
<view-3>  
 <employee\_name>{$e/fname}{$e/lname}</employee\_name>  
 {  
 for $d in doc("Q3-company.xml")/companyDB/departments/department[@dno=$e/@worksFor]  
 where $d/dname="Research"  
 return $d/dname  
 }  
 {  
 for $s in doc("Q3-company.xml")/companyDB/employees/employee[@ssn=$e/@supervisor]  
 return   
 <supervisor>{$s/fname}{$s/lname}</supervisor>  
 }  
 {  
 $e/salary  
 }  
</view-3>

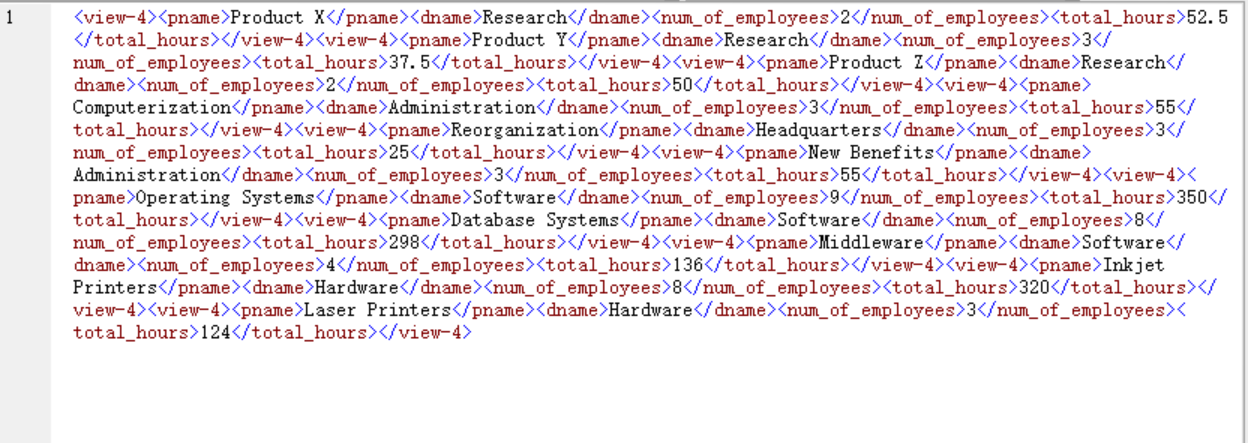
**

**

1. A view that has the project name, controlling department name, number of employees, and total hours worked per week on the project for each project.

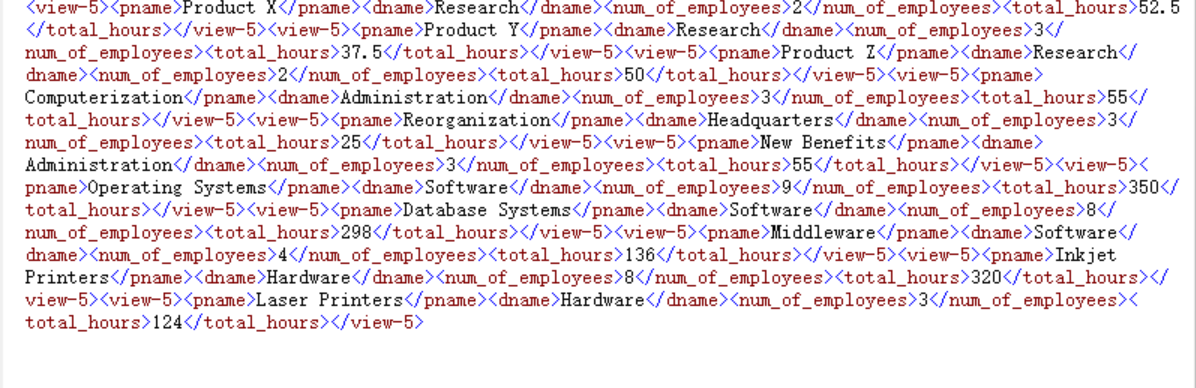
for $p in doc("Q3-company.xml")/companyDB/projects/project  
return  
 <view-4>  
 {$p/pname}  
 {  
 for $d in doc("Q3-company.xml")/companyDB/departments/department[@dno=$p/@controllingDepartment]  
 return $d/dname  
 }  
 <num\_of\_employees>  
 {  
 fn:count(  
 for $e in doc("Q3-company.xml")/companyDB/employees/employee[@ssn=$p/workers/worker/@essn]  
 return $e  
 )}  
 </num\_of\_employees>  
 <total\_hours>  
 {  
 fn:sum(  
 for $w in doc("Q3-company.xml")/companyDB/projects/project[pname=$p/pname]  
 return $w/workers/worker  
 )}  
 </total\_hours>  
 </view-4>





1. A view that has the project name, controlling department name, number of employees, and total hours worked per week on the project for each project with more than one employee working on it.
2. for $p in doc("Q3-company.xml")/companyDB/projects/project  
   return  
    <view-5>  
    {$p/pname}  
    {  
    for $d in doc("Q3-company.xml")/companyDB/departments/department[@dno=$p/@controllingDepartment]  
    return $d/dname  
    }  
    <num\_of\_employees>  
    {  
    fn:count(  
    for $e in doc("Q3-company.xml")/companyDB/employees/employee[@ssn=$p/workers/worker/@essn]  
    return $e  
    )}  
    </num\_of\_employees>  
    <total\_hours>  
    {  
    fn:sum(  
    for $w in doc("Q3-company.xml")/companyDB/projects/project[pname=$p/pname]  
    return $w/workers/worker  
    )}  
    </total\_hours>  
    </view-5>





**4: Object-Relational Database (10 Marks)**

Answer questions **4(a), 4(b), and 4(c)** by considering the following UML diagram for the online course application we have used in the lab for Object-Relational Mapping.

Diagram

Description automatically generated

**Q4(a)** Note that Instructor and learner have Many-To-Many relationships with the course. We created class models for User, Instructor, Learner, Course, and Lesson to directly map these entities to objects. In addition, we created an Enrollment class to allow students to get enrolled for a course. We have added Instructors in the Course class in the lab. In this task, complete the below code by updating the Course class by adding Many-To-Many relationship with Learner via Enrollment relationship. **(2 marks)**

**# Update Course model**

class Course(models.Model):

name = models.CharField(null=False, max\_length=100, default='online course')

description = models.CharField(max\_length=500)

# Many-To-Many relationship with Instructors

instructors = models.ManyToManyField(Instructor)

*# Complete below code for Many-To-Many relationship with Learner via Enrollment relationship, similar to instructors above, however you have to lookup learner via enrollment class.*

learners = *models.ManyToManyField(Learner)*

def \_\_str\_\_(self):

return "Name: " + self.name + "," + \

"Description: " + self.description

**Q4(b) A**fter defining class models in the lab assignment, we created objects by defining methods for User, Instructor, Learner, Course, and Lesson using the ‘write.py’ file. We also defined a method for creating objects for Course\_Enrollment relationship. In this task, you are required to open the ‘write.py’ file and append a method for Course\_instructor\_relationship() to add courses for instructors. You are required to complete the below code. **(3 marks)**

def write\_course\_instructor\_relationships():

# Get at least three instructors [Hint: use objects.get() method on the Instructor class

# Your code should go here

instructor\_user\_peter = Instructor.objects.get(first\_name='Peter')  
instructor\_user\_yan=Instructor.objects.get(first\_name='Yan')  
instructor\_user\_joy=Instructor.objects.get(first\_name='Joy')

# Get at least two related courses [Hint: use objects.get() method on the Course class

# Your code should go here

course\_cloud\_app = Course.objects.get(name\_\_contains='Cloud')  
course\_python = Course.objects.get(name\_\_contains='Python')

# Add instructors to courses [hint: use Add() method on the Instructor class]

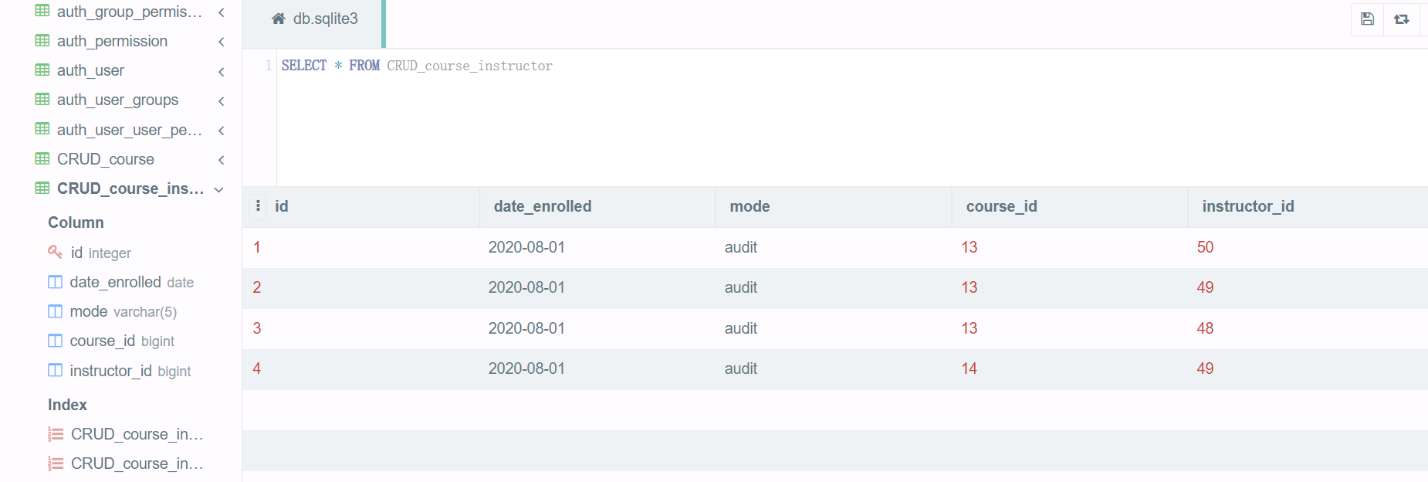
# Your code should go here

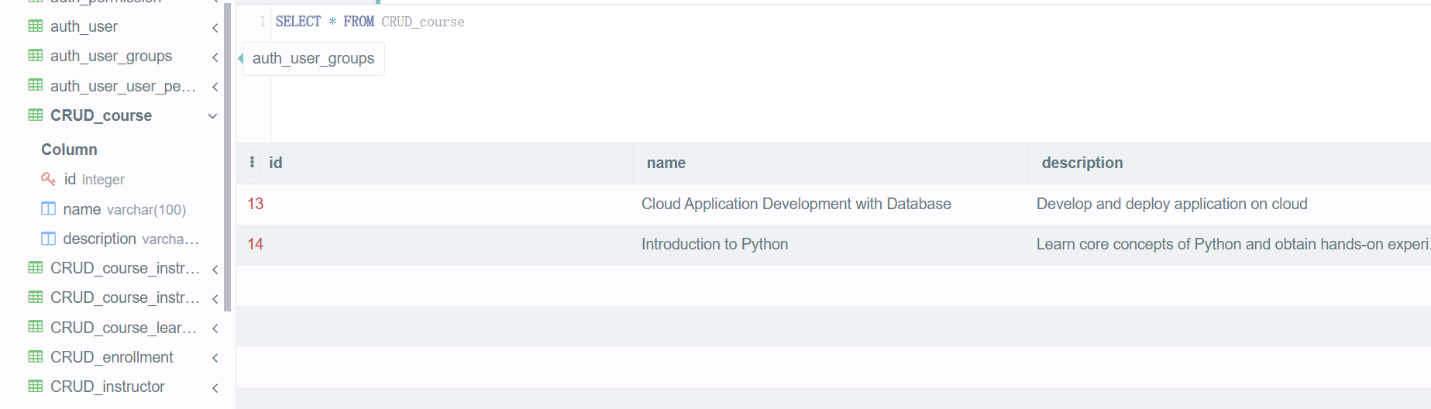
john\_cloud=Course\_Instructor.objects.create(instructor=instructor\_user\_peter, date\_enrolled=date(2020, 8, 1),  
 course=course\_cloud\_app, mode='audit'  
 )  
john\_cloud.save()  
  
joy\_cloud=Course\_Instructor.objects.create(instructor=instructor\_user\_joy, date\_enrolled=date(2020, 8, 1),  
 course=course\_cloud\_app, mode='audit'  
 )  
joy\_cloud.save()  
  
yan\_cloud = Course\_Instructor.objects.create(instructor=instructor\_user\_yan, date\_enrolled=date(2020, 8, 1),  
 course=course\_cloud\_app, mode='audit'  
 )  
yan\_cloud.save()  
  
joy\_python = Course\_Instructor.objects.create(instructor=instructor\_user\_joy, date\_enrolled=date(2020, 8, 1),  
 course=course\_python, mode='audit'  
 )  
joy\_python.save()  
  
print("Course-instructor relationships saved... ")

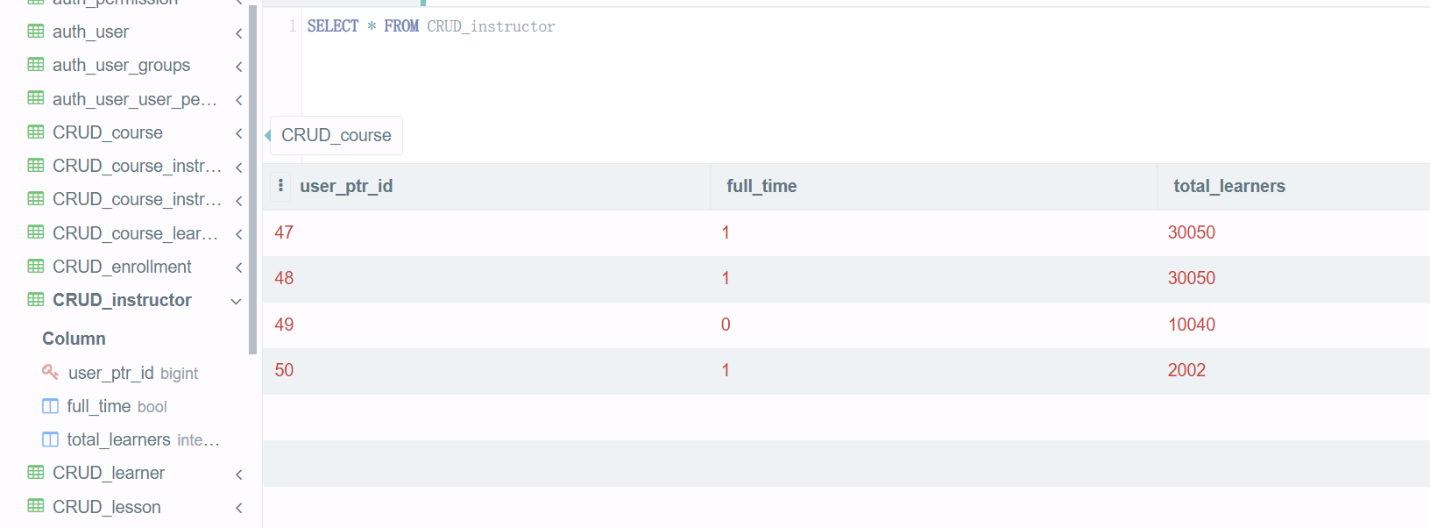
write\_course\_instructor\_relationships()

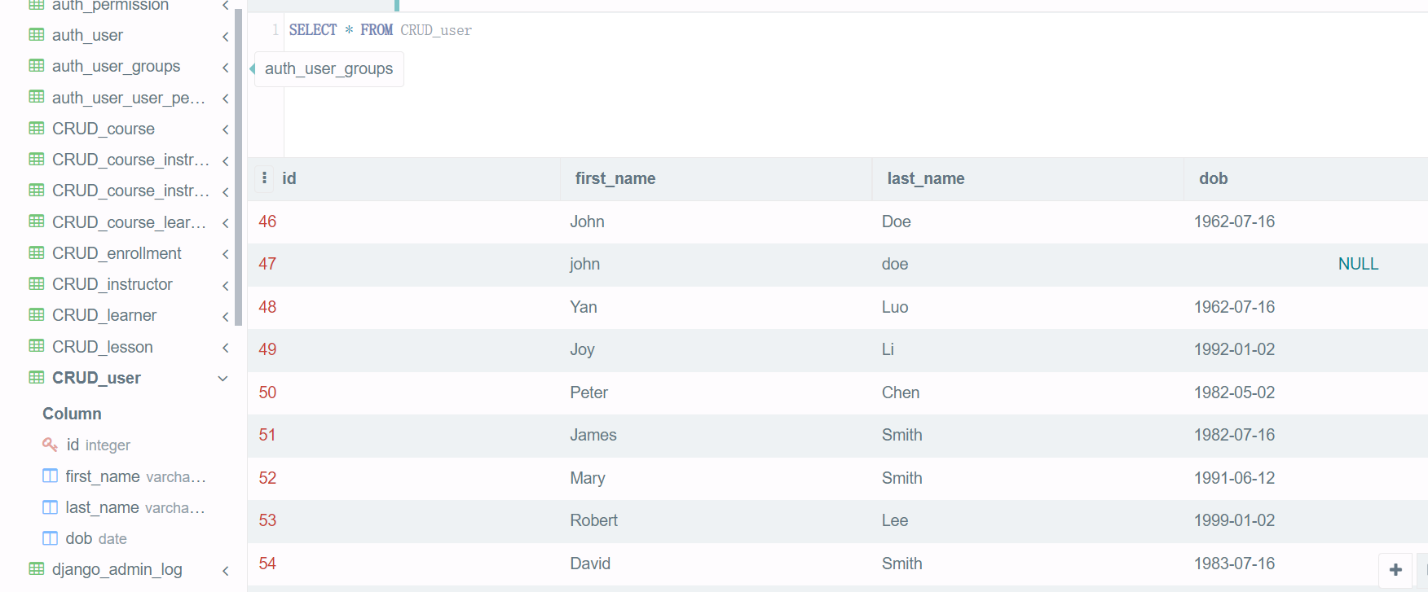
Save the ‘write.py’ file and run it in the terminal to populate the ‘course\_instructor’ table in the database. Provide a screenshot of course\_instructor table with three rows referring to rows in the course and instructor tables.

Save the ‘write.py’ file and run it in the terminal to populate the ‘course\_instructor’ table in the database. Provide a screenshot of course\_instructor table with three rows referring to rows in the course and instructor table. Also, provide a screenshot of the course and user table to refer to rows in these tables corresponding to the given instructor of the course from the course\_instructor table.









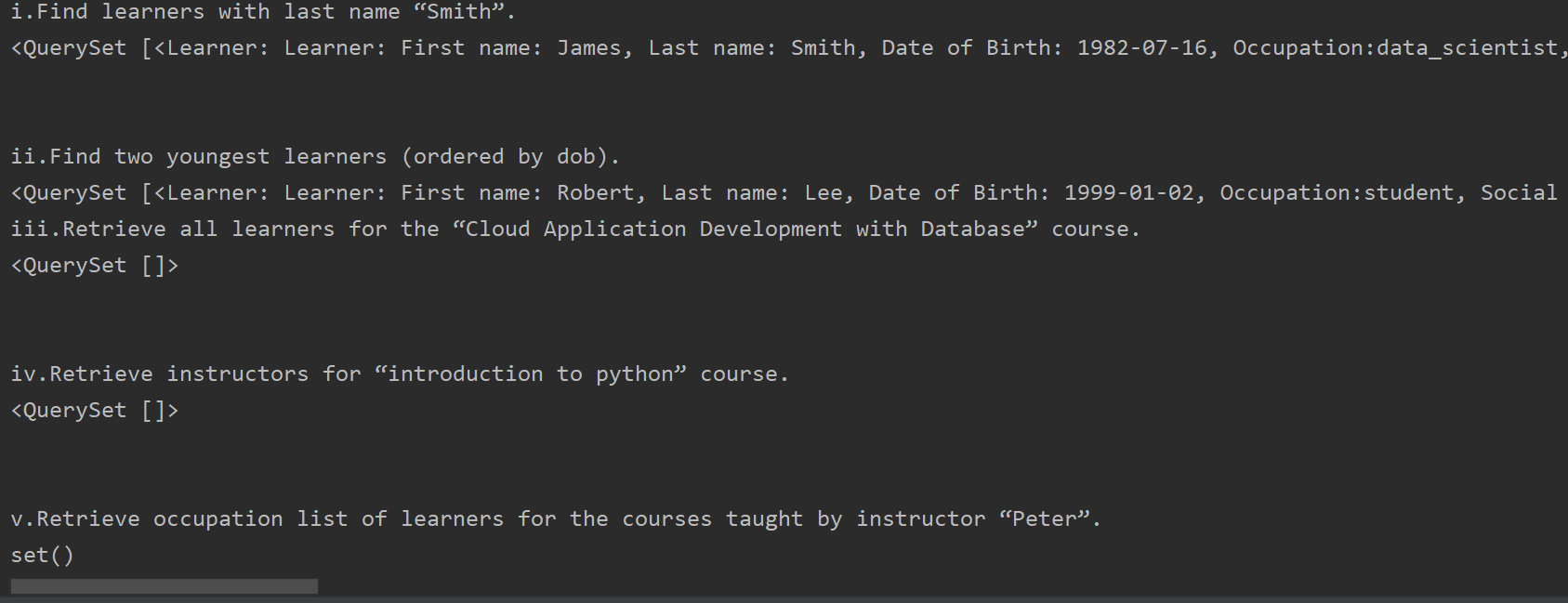
**Q4(c)** At this point, you have defined models and created objects. Now you will apply your skills of querying objects. In this task, you are required to create ‘Read.py’ file under your ‘online\_Courses’ project directory.

Open the ‘Read.py’ file; first, copy Django-specific settings from the ‘wsgi.py’ file, set database connection, and import model classes, as we did in the lab session. Create code to perform the following queries: **(1\*5 = 5 marks)**

1. Find learners with last name “Smith”.
2. Find two youngest learners (ordered by dob).
3. Retrieve all learners for the “Cloud Application Development with Database” course.
4. Retrieve instructors for “introduction to python” course.
5. Retrieve occupation list of learners for the courses taught by instructor “Peter”.

Format your code in a way that the query results should look like as given below. Provide a screenshot of your code and query results.

# Django specific settings  
import inspect  
import os  
os.environ.setdefault('DJANGO\_SETTINGS\_MODULE', 'OnlineCourse.settings')  
from django.db import connection  
# Ensure settings are read  
from django.core.wsgi import get\_wsgi\_application  
application = get\_wsgi\_application()  
  
from CRUD.models import \*  
from datetime import date  
  
  
print('i.Find learners with last name “Smith”.')  
learners\_smith = Learner.objects.filter(last\_name='Smith')  
print(learners\_smith)  
print("\n")  
  
print('ii.Find two youngest learners (ordered by dob).')  
learners = Learner.objects.order\_by('-dob')[0:2]  
print(learners)  
  
  
print('iii.Retrieve all learners for the “Cloud Application Development with Database” course.')  
cloud\_learners=Learner.objects.filter(course\_\_name\_\_contains='Cloud')  
print(cloud\_learners)  
print("\n")  
  
print('iv.Retrieve instructors for “introduction to python” course.')  
introduction\_to\_python\_instructors=Instructor.objects.filter(course\_\_name\_\_contains='Python')  
print(introduction\_to\_python\_instructors)  
print("\n")  
  
print('v.Retrieve occupation list of learners for the courses taught by instructor “Peter”.')  
courses=Course.objects.filter(instructors\_\_first\_name='Peter')  
occupation\_list = set()  
for course in courses:  
 for learner in course.learners.all():  
 occupation\_list.add(learner.occupation)  
print(occupation\_list)  
print("\n")



1. Find learners with last name `Smith`:

<QuerySet [

<Learner: First name: James, Last name: Smith, Date of Birth: 1982-07-16, Occupation: data\_scientist, Social Link: https://www.linkedin.com/james/>,

<Learner: First name: Mary, Last name: Smith, Date of Birth: 1991-06-12, Occupation: dba, Social Link: <https://www.facebook.com/mary/>>,

<Learner: First name: David, Last name: Smith, Date of Birth: 1983-07-16, Occupation: developer, Social Link: https://www.linkedin.com/david/>,

<Learner: First name: John, Last name: Smith, Date of Birth: 1986-03-16, Occupation: developer, Social Link: <https://www.linkedin.com/john/>>]>

2. Find top two youngest learners:

<QuerySet [

<Learner: First name: Robert, Last name: Lee, Date of Birth: 1999-01-02, Occupation: student, Social Link: <https://www.facebook.com/robert/>>,

<Learner: First name: Mary, Last name: Smith, Date of Birth: 1991-06-12, Occupation: dba, Social Link: <https://www.facebook.com/mary/>>]

3. Retrieve all learners for “Cloud Application Development with Database" course

<QuerySet [

>/

Retrieve instructors for “introduction to python" course

<QuerySet [

>/

4. Retrieve instructors for “introduction to python" course

<QuerySet [

>/

5. Retrieve occupation list of learners for the courses taught by instructor “Peter”

<QuerySet [

>/

**5: Data Warehousing and OLAP (10 Marks)**

**Q5** In this task, you will perform “Online Analytical Processing” (OLAP) style queries over a simple “star schema” given below. The corresponding schema, along with the database, is provided with this assignment. You are required to load the schema in Oracle Database and use SQL developer to perform the following queries: **(2\*5=10 marks)**



1. Summarize (sum, min, and count) store sales for USA and Canada in 2016 by store zip code and month without using subtotal operators. Provide query code and results.

select STOREZIP, TIMEMONTH, sum(SALESDOLLAR), min(SALESDOLLAR), count(SALESNO)

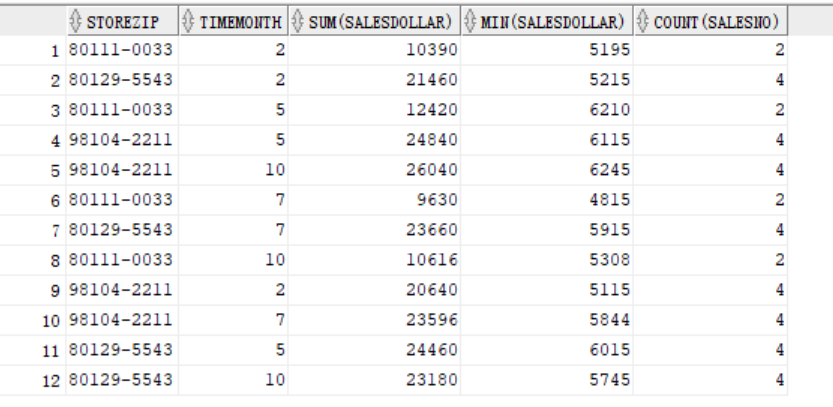
from sssales join ssstore on sssales.storeid=ssstore.storeid

join sstimedim on sssales.TIMENO=sstimedim.TIMENO

where STORENATION in ('USA', 'Canada')

and sstimedim.TIMEYEAR=2016

group by STOREZIP, TIMEMONTH;



1. Use the same problem in part one and generate all possible subtotals by zip code and month using the data CUBE operator. Provide query code and results in your report.

select STOREZIP, TIMEMONTH, sum(SALESDOLLAR), min(SALESDOLLAR), count(SALESNO)

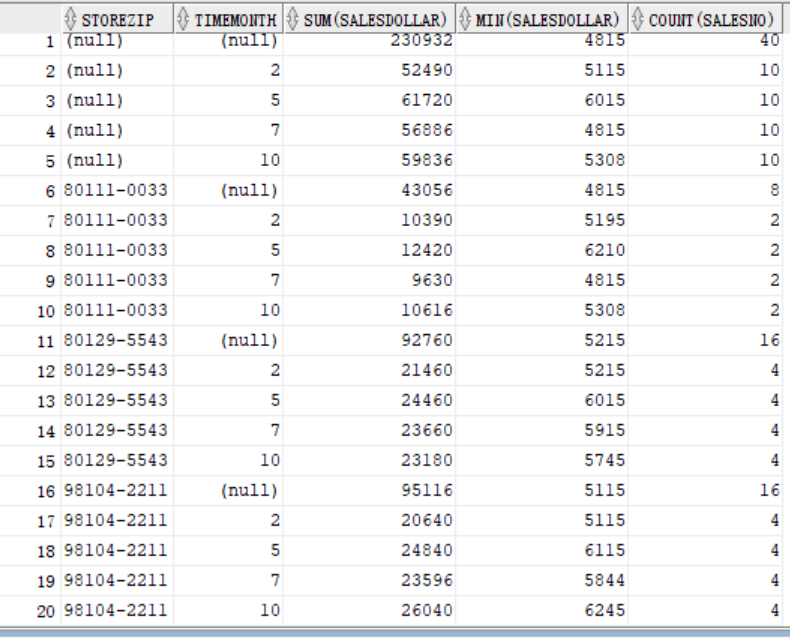
from sssales join ssstore on sssales.storeid=ssstore.storeid

join sstimedim on sssales.TIMENO=sstimedim.TIMENO

where STORENATION in ('USA', 'Canada')

and sstimedim.TIMEYEAR=2016

group by cube(STOREZIP, TIMEMONTH);



1. Summarize (SUM, COUNT, and MIN) store sales for USA and Canada between 2016 and 2017 by year and month. Generate partial subtotals for *year* and *month* using the ROLLUP operator. Provide query code and results in your report.

select TIMEYEAR, TIMEMONTH, sum(SALESDOLLAR), min(SALESDOLLAR), count(SALESNO)

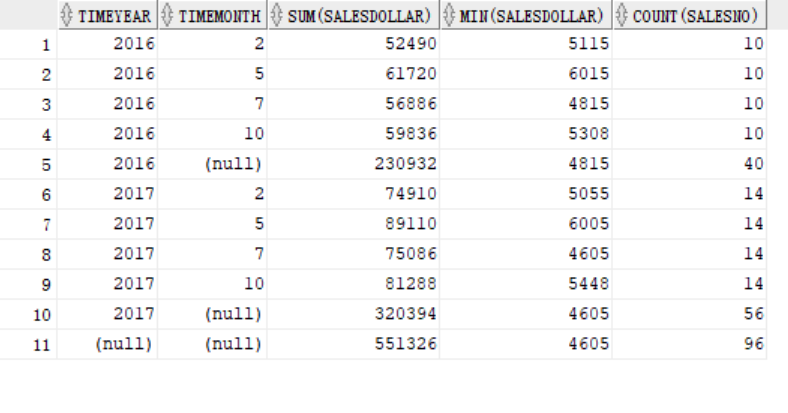
from sssales join ssstore on sssales.storeid=ssstore.storeid

join sstimedim on sssales.TIMENO=sstimedim.TIMENO

where STORENATION in ('USA', 'Canada')

and sstimedim.TIMEYEAR between 2016 and 2017

group by rollup(TIMEYEAR, TIMEMONTH);



1. Rewrite the query in part 3, and provide a set of subtotals on year, quarter, and month. Sort the results in a convenient order.

select TIMEYEAR,TIMEQUARTER, TIMEMONTH, sum(SALESDOLLAR), min(SALESDOLLAR), count(SALESNO)

from sssales join ssstore on sssales.storeid=ssstore.storeid

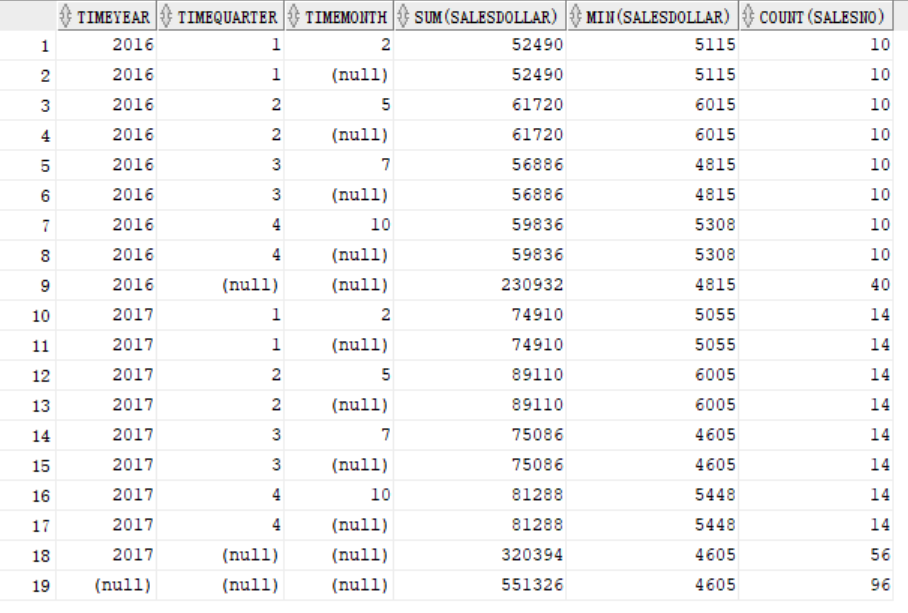
join sstimedim on sssales.TIMENO=sstimedim.TIMENO

where STORENATION in ('USA', 'Canada')

and sstimedim.TIMEYEAR between 2016 and 2017

group by rollup(TIMEYEAR,TIMEQUARTER, TIMEMONTH)

order by TIMEYEAR,TIMEQUARTER, TIMEMONTH;



1. Identify differences between results containing the GROUP BY clause (but not subtotal operators) and data cube and Rollup Results from queries in 1, 2, and 3. You should compare and contrast results based on the number of rows and subtotals groups. Explain why or why not each type of subtotal is useful for supporting business-specific queries.

Group by can not get subtotals of the query result, and cube can get all possible subtotals of the query result, rollup can get partial subtotals.

When we need to get all subtotals, cube is useful for supporting it. Such as queries in 2, there’re two different dimensions zipcode and month, it’s better use cube.

And when we just need some subtotals such as queries in 4(year, quarter, month), they all in a dimension, we can use rollup.

---------------------------------------------------**End of Assessment Task 1**-----------------------------------------------